

<p><b>Rules to Answer</b></p> <p>Accurate answer supported by evidence, and reasoning links answer and evidence. Response is specific and detailed. Presence of keywords and scientific terminology.</p>	<p><b>Rules to Answer</b></p> <p>Refer content if unable to recall after 15-30sec of trying, then again retrieve from memory without help.</p>	<p><b>Primitive Data Types</b></p> <p>Free Recall</p>	<p><b>Primitive Data Types</b></p> <p>Data types that are defined by the system. e.g. int, float, double, bool</p>
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<p><b>Asymptotic Notation</b></p> <p>Elaborate</p>	<p><b>Asymptotic Notation or Algorithm's growth rate</b></p> <p>are languages that allow us to analyze an algorithm's running time by identifying its behavior as the input size for the algorithm increases.</p>	<p><b>Big-O (O) vs Big-Omega (Ω)</b></p> <p>Main idea</p>	<p><b>Big-O</b></p> <p>Asymptotic Notation for the worst case, or ceiling of growth for a given function. Analogy: At most</p>
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<p><b>Big-O</b></p> <p>Provides us with an <b>asymptotic upper bound</b> for the growth rate of runtime of an algorithm</p>	<p><b>Big-O</b></p> <p>Say <math>f(n)</math> is your algorithm runtime, and <math>g(n)</math> is an arbitrary time complexity you are trying to relate to your algorithm</p>	<p><b>Big-O</b></p> <p><math>f(n)</math> is <math>O(g(n))</math>, if for some constants <math>c</math> (<math>c &gt; 0</math>) and <math>n_0</math>, <math>f(n) \leq c * g(n)</math> for every input size <math>n(n &gt; n_0)</math></p>	<p><b>Big-O Examples</b></p>
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<p><b>Big-O Examples</b></p> <p><math>f(n) = 3 \log n + 100</math> <math>g(n) = \log n</math></p>	<p><b>Big-Omega</b></p> <p>Best case Floor growth rate Asymptotic lower bound Analogy: At least</p>	<p><b>Big O and Big-Omega</b></p> <p>Give the mathematical definition</p>	<p><b>Small-o and Small-omega (ω)</b></p> <p>Elaborate</p>
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<p><b>Small-o and Small-omega</b></p> <p>Both are not asymptotically tight (no equal to sign)</p>	<p><b>Theta Θ</b></p> <p>Explain</p>	<p><b>Theta</b></p> <p>Asymptotical tight bound</p> <p><math>c1 * g(n) &lt; f(n) &lt; c2 * g(n)</math></p>	<p><b>Recurrence</b></p> <p>Main idea</p>
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<p><b>Recurrence</b></p> <p>An equation that recursively defines a sequence once one or more initial terms are given: each further term of the sequence or array is defined as a function of the preceding terms.</p>	<p><b>Recursion vs Recurrence</b></p> <p>Explain</p>	<p><b>Recursion vs Recurrence</b></p> <p>A recurrence relation uses recursion to create a sequence. Recursion is not limited to generation of sequences.</p>	<p><b>Recursion</b></p> <p>Recursion is the repeated use of a procedure or action</p>
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<p><b>Recurrence</b></p> <p>In fact, a recurrence relation uses recursion to define a sequence.</p>	<p><b>Recurrence</b></p> <p>This sequence is built in such a way that each term is defined as a combination of previous terms. The generation of such a sequence is a requirement in the definition</p>	<p><b>Recurrence</b></p> <p>Recurrence is analyzed by telescoping</p>	<p><b>Binary Search</b></p> <p>Main Idea</p>
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<p><b>Binary Search</b></p> <p>Search algorithm that finds the position of a target value within a sorted array</p>	<p><b>Binary Search</b></p> <p>Elements are already sorted in ascending order Divide and Conquer algorithm is used</p>	<p><b>Binary Search</b></p> <p>Element <math>x</math> is either in the left half of the array or in the right half or not there at all</p>	<p><b>Binary Search</b></p> <p>Compare <math>x</math> to middle element <math>k</math>, <math>x &gt; k</math> (<math>x</math> not in left half), <math>x &lt; k</math> (<math>x</math> not in right half)</p>
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### Running time of Binary Search

Explain

### Running time of Binary Search

$\lg_2 N$   
binary search runs in  $O(\lg N)$

### One word answer

In computer science you use  $\lg_2 N$  or  $\lg_{10} N$

In computer science you use  $\lg_2 N$  or  $\lg_{10} N$

$\lg_2 N$

### Merging

Elaborate

### Merging

Merging is not a divide and conquer algorithm, but part of mergesort algorithm

### Merge Sort

Merge Sort is a divide and conquer algorithm

### Merging

Take two sorted arrays of numbers and make a single array which is sorted of all of those numbers

### Merge Sort

Type of Algorithm

### Merge Sort

Explain the complete algorithm

### Mergesort

First divide the list into smallest unit (1 element), then compare each element with the adjacent list to sort and merge two adjacent lists. Finally all the elements are sorted and merged.

### Run time for mergesort

### Run time for mergesort

$O(n \cdot \lg n)$

### Insertion sort vs Mergesort

Which one is better and why?

### Insertion sort vs Mergesort

Insertion sort is  $O(n^2)$ .  
Mergesort is  $O(n \cdot \lg n)$ .  
So, merge-sort is a superior algorithm in terms of its running time on large datasets.

### Algorithm runtime

Insertion sort and Bubble sort

### Insertion and bubble sort Runtime

Both has  
Best:  $\Omega(n)$   
Worst:  $O(n^2)$